

Journal of Statistics & Management Systems		
Transmitting of		personal care
_		
1.		terter in
11.1000-012	the last set of	
10000	1.0.0	0.005. 144
121022-000		100 million (100
1.1.1		
19+ *		@3557**

Journal of Statistics and Management Systems

ISSN: (Print) (Online) Journal homepage: https://www.tandfonline.com/loi/tsms20

# E-duality results for E-differentiable vector optimization problems involving E-type I functions

Najeeb Abdulaleem

To cite this article: Najeeb Abdulaleem (2021): E-duality results for E-differentiable vector optimization problems involving E-type I functions, Journal of Statistics and Management Systems, DOI: 10.1080/09720510.2021.1890907

To link to this article: https://doi.org/10.1080/09720510.2021.1890907



Published online: 23 Jul 2021.



🕼 Submit your article to this journal 🗗

Article views: 11



View related articles



🌔 🛛 View Crossmark data 🗹

*Journal of Statistics & Management Systems* ISSN 0972-0510 (Print), ISSN 2169-0014 (Online) DOI : 10.1080/09720510.2021.1890907



# *E*-duality results for *E*-differentiable vector optimization problems involving *E*-type I functions

Najeeb Abdulaleem Department of Mathematics Hadhramout University P. O. Box (50511-50512) Al-Mahrah Yemen

and

Faculty of Mathematics and Computer Science University of Łódź Banacha 22, 90–238 Łódź Poland

#### Abstract

In this paper, the class of *E*-differentiable vector optimization problems is considered. The so-called vector Wolfe and mixed *E*-duality problems are defined for the considered *E*-differentiable multiobjective programming problem and several Wolfe and mixed *E*-duality theorems are derived under appropriate *E*-type I functions.

Subject Classification: 90C26, 90C30, 26B25.

*Keywords:* E-type I functions, E-differentiable function, E-optimality conditions, Wolfe E-duality, Mixed E-duality.

# 1. Introduction

Consider the following multiobjective programming problem with inequality constraints:

E-mail: nabbas985@gmail.com



minimize 
$$f(x) = (f_1(x), ..., f_p(x))$$
  
subject to  $g_t(x) \leq 0, t \in T = \{1, ..., m\},$  (VP)

where  $f_i : \mathbb{R}^n \to \mathbb{R}$ ,  $i \in I = \{1, ..., p\}$ ,  $g_t : \mathbb{R}^n \to \mathbb{R}$ ,  $t \in T$ , are (not necessarily) differentiable functions on  $\mathbb{R}^n$ . Let *D* denote the set of all feasible solutions in problem (VP), that is,

$$D := \{x \in R^n : g_t(x) \leq 0, t \in T\}$$

Further, we denote by  $T(\overline{x}) = \{t \in T : g_t(\overline{x}) = 0\}$  the set of indices of inequality constraints that are active at the arbitrary feasible point  $\overline{x}$ .

The following convention for equalities and inequalities will be used in the paper. If  $x = (x_1, x_2, ..., x_n)^T$  and  $y = (y_1, y_2, ..., y_n)^T$  in  $\mathbb{R}^n$ , then,  $x = y \Leftrightarrow x_i = y_i$ , i = 1, 2, ..., n;  $x > y \Leftrightarrow x_i > y_i$ , i = 1, 2, ..., n;  $x \ge y \Leftrightarrow x_i \ge y_i$ , i = 1, 2, ..., n;  $x \ge y \Leftrightarrow x_i \ge y_i$ , i = 1, 2, ..., n;  $x \ge y \Leftrightarrow x \ge y$  and  $x \ne y$ .

For the multiobjective programming problem (VP), the following concepts of a weakly efficient solution (a weak Pareto solution) and an efficient solution (a Pareto solution) are defined as follows:

**Definition 1.1 : [1]** A point  $\overline{x}$  is said to be a weakly efficient solution (weak Pareto solution) of (VP) if there exists no *x* such that

 $f(x) < f(\overline{x}).$ 

**Definition 1.2 : [1]** A point  $\overline{x}$  is said to be an efficient solution (a Pareto solution) of (VP) if there exists no *x* such that

 $f(x) \le f(\overline{x}).$ 

During the last few years, there has been an increasing interest in generalizations of convexity in connection with sufficiency and duality in optimization problems. Optimality conditions and duality results for differentiable and nondifferentiable multiobjective programming problems have been studied extensively in the literature (see, for example, [1], [2], [4], [5], [6], [8], [9], [11], [12], [13], [14], [15] and others).

Hanson and Mond [10] introduced type I and type II invexities which have been generalized to pseudo type I and quasi type I functions by Reuda and Hanson [13]. Later, Aghezzaf and Hachimi [8] introduced vector-valued generalized type I functions. Megahed et al. [12] introduced a new definition of an *E*-differentiable convex function which transforms a (not necessarily) differentiable function to a differentiable function under an operator  $E: \mathbb{R}^n \to \mathbb{R}^n$ . Abdulaleem [3] introduced the *E*-differentiable *E*-invexity notion as a generalization of the concept of *E*-differentiable *E*-convexity. Recently, Abdulaleem [7] introduced a new concept of the so-called *E*-type I functions for *E*-differentiable vector optimization problems. Namely, Abdulaleem defined several classes of generalized vector-valued *E*-type I functions for *E*-differentiable functions as a generalization of the concept of differentiable vector-valued type I functions introduced by Hanson and Mond [10].

In this paper, we consider a new class of (not necessarily differentiable) multiobjective programming problems in which the involved functions are *E*-differentiable. For the considered *E*-differentiable programming problem, its vector *E*-dual problems in the sense of Wolfe and mixed are defined. Then various Wolfe and mixed *E*-duality theorems are established between the considered *E*-differentiable multicriteria optimization problem and its vector Wolfe and mixed *E*-dual problems under *E*-type I functions.

## 2. Preliminaries

We now give the definition of an *E*-invex set introduced by Abdulaleem [3].

**Definition 2.1: [3]** Let  $E: \mathbb{R}^n \to \mathbb{R}^n$ . A set  $M \subseteq \mathbb{R}^n$  is said to be an *E*-invex set if and only if there exists a vector-valued function  $\eta: M \times M \to \mathbb{R}^n$  such that the relation

$$E(u) + \lambda \eta (E(x), E(u)) \in M$$

holds for all  $x, u \in M$  and any  $\lambda \in [0, 1]$ .

We now recall for a common reader the definition of an *E*-differentiable function introduced by Megahed et al. [12].

**Definition 2.2:** [12] Let  $E: \mathbb{R}^n \to \mathbb{R}^n$  and  $f: M \to \mathbb{R}$  be a (not necessarily) differentiable function at a given point  $u \in M$ . It is said that f is an *E*-differentiable function at u if and only if  $f \circ E$  is a differentiable function at u (in the usual sense), that is,

$$(f \circ E)(x) = (f \circ E)(u) + \nabla (f \circ E)(u)(x-u) + \theta(u, x-u) ||x-u||,$$
(1)

where  $\theta(u, x - u) \rightarrow 0$  as  $x \rightarrow u$ .

Let  $E: \mathbb{R}^n \to \mathbb{R}^n$  be a given one-to-one and onto operator. Throughout the paper, we shall assume that the functions constituting the problem (VP) are *E*-differentiable at any feasible solution. Now, for the problem

(VP), we define its associated differentiable vector optimization problem  $(VP_r)$  as follows:

minimize 
$$f(E(x)) = (f_1(E(x)), ..., f_p(E(x)))$$
  
subject to  $g_t(E(x)) \le 0, t \in T = \{1, ..., m\}.$  (VP<sub>E</sub>)

Let  $D_E$  denote the set of all feasible solutions in problem (VP<sub>E</sub>), that is,

$$D_E := \{x \in R^n : g_t(E(x)) \leq 0, t \in T\}.$$

Further, we denote by  $T_E(\overline{x}) = \{t \in T : g_t(E(\overline{x})) = 0\}$  the set of indices of inequality constraints that are active at the arbitrary feasible point  $\overline{x} \in D_E$ .

**Definition 2.3 :** [1] A point  $E(\overline{x}) \in D$  is said to be a weakly *E*-efficient solution (weak *E*-Pareto solution) of (VP) if there exists no  $E(x) \in D$  such that

$$f(E(x)) < f(E(\overline{x})).$$

**Definition 2.4 : [1]** A point  $E(\overline{x}) \in D$  is said to be an *E*-efficient solution (an *E*-Pareto solution) of (VP) if there exists no  $E(x) \in D$  such that

$$f(E(x)) \le f(E(\overline{x})).$$

**Lemma 2.5 :** [1] Let  $E: \mathbb{R}^n \to \mathbb{R}^n$  be a one-to-one and onto. Then  $E(D_E) = D$ .

**Lemma 2.6 :** [1] Let  $\overline{x} \in D$  be a weakly efficient solution (an efficient solution) of the problem (VP). Then, there exists  $\overline{z} \in D_E$  such that  $\overline{x} = E(\overline{z})$  and  $\overline{z}$  is a weakly efficient solution (an efficient solution) of the problem (VP<sub>E</sub>).

**Lemma 2.7:** [1] Let  $\overline{z} \in D_E$  be a weakly efficient solution (an efficient solution) of the problem  $(VP_E)$ . Then  $E(\overline{z})$  is a weakly efficient solution (an efficient solution) of the problem (VP).

**Definition 2.8:** [7] Let  $E: \mathbb{R}^n \to \mathbb{R}^n$ ,  $M \subseteq \mathbb{R}^n$  be an open *E*-invex set. (*f*, *g*) is said to be (strictly) *E*-type I with respect to  $\eta$  at  $u \in M$  if there exists a vector-valued function  $\eta: M \times M \to \mathbb{R}^n$  such that, for all  $x \in M$ ,

$$f_i(E(x)) - f_i(E(u)) \ge \nabla f_i(E(u))\eta(E(x), E(u)), \quad (>)$$

$$(2)$$

$$-g_t(E(u)) \ge \nabla g_t(E(u))\eta(E(x), E(u)).$$
(3)

**Theorem 2.9 : [1]** (*E*-Karush-Kuhn-Tucker necessary optimality conditions). Let  $\overline{x} \in D_E$  be a weakly efficient solution of the problem  $(VP_E)$  (and, thus,  $E(\overline{x})$  be a weakly *E*-efficient solution of the considered multiobjective programming problem (VP)). Further, f and g are *E*-differentiable at  $\overline{x}$  and the Guignard constraint qualification [3] is satisfied at  $\overline{x}$ . Then there exist Lagrange multipliers  $\overline{\tau} \in \mathbb{R}^p$ ,  $\overline{\xi} \in \mathbb{R}^m$  such that

$$\sum_{i=1}^{p} \overline{\tau}_{i} \nabla f_{i} \left( E(\overline{x}) \right) + \sum_{t=1}^{m} \overline{\xi}_{t} \nabla g_{t} \left( E(\overline{x}) \right) = 0, \qquad (4)$$

$$\xi_t g_t \left( E(\overline{x}) \right) = 0, t \in T(E(\overline{x})), \tag{5}$$

$$\overline{\tau} \geq 0, \, \overline{\xi} \geq 0. \tag{6}$$

#### 3. Wolfe *E*-duality

In this section, for the considered *E*-differentiable multiobjective programming problem (VP), we define the following dual problem in the sense of Wolfe. Further, under *E*-type I functions, we prove *E*-duality results between the multiobjective programming problem (VP) and its vector *E*-dual problem in the sense of Wolfe. Let  $E: \mathbb{R}^n \to \mathbb{R}^n$  be a one to one and onto operator. We define for problem (VP<sub>*E*</sub>) its *E* -differentiable vector *E*-dual problem in the sense of Wolfe as follows:

$$\begin{aligned} \text{maximize } f(E(y)) + \left[\sum_{t=1}^{m} \xi_{t} g_{t}(E(y))\right] e \\ \text{subject to } \sum_{i=1}^{p} \tau_{i} \nabla(f_{i} \circ E)(y) + \sum_{t=1}^{m} \xi_{t} \nabla(g_{t} \circ E)(y) = 0, \qquad (\text{WD}_{E}) \\ y \in \mathbb{R}^{n}, \tau \geq 0, \tau \in \mathbb{R}^{p}, \tau e = 1, e = (1, 1, ..., 1)^{T} \in \mathbb{R}^{p}, \\ \xi \in \mathbb{R}^{m}, \xi \geq 0, \end{aligned}$$

where  $e = (1, ..., 1) \in \mathbb{R}^p$ . Denote by

$$\Gamma_{E} = \left\{ (y, \tau, \xi) \in \mathbb{R}^{n} \times \mathbb{R}^{p} \times \mathbb{R}^{m} : \sum_{i=1}^{p} \tau_{i} \nabla (f_{i} \circ E)(y) + \sum_{t=1}^{m} \xi_{t} \nabla (g_{t} \circ E)(y) = 0, \tau e = 1, \tau \ge 0, \xi \ge 0 \right\}.$$

the feasible solutions set of  $(WD_E)$ . Further,  $Y_E = \{y \in \mathbb{R}^n : (y, \tau, \xi) \in \Gamma_E\}$ .

**Theorem 3.1 :** (Wolfe weak duality between  $(VP_E)$  and  $(WD_E)$ ). Let  $x \in D_E$  and  $(y, \tau, \xi) \in \Gamma_E$  such that (f, g) is E-type I at y. Then

$$f(E(x)) \not< f(E(y)) + \left[\sum_{t=1}^{m} \xi_t g_t(E(y))\right] e.$$
(7)

**Proof** : Let  $x \in D_E$  and  $(y, \tau, \xi) \in \Gamma_E$ . We proceed by contradiction. Suppose, contrary to the result, that the inequality

$$f(E(x)) < f(E(y)) + \left[\sum_{t=1}^{m} \xi_t g_t(E(y))\right] e$$
 (8)

holds. Thus,

$$f(E(x)) < f(E(y)) + \left[\sum_{t=1}^{m} \xi_{t} g_{t}(E(y))\right] \sum_{i=1}^{p} \tau_{i} , i \in I$$
(9)

holds. Multiplying the above inequality by  $\tau_i$  with taking that  $\sum_{i=1}^{p} \tau_i = 1$ , the following inequality

$$\sum_{i=1}^{p} \tau_{i} f_{i}(E(x)) < \sum_{i=1}^{p} \tau_{i} f_{i}(E(y)) + \sum_{t=1}^{m} \xi_{t} g_{t}(E(y))$$
(10)

holds. From  $x \in D_E$ , it follows that

$$\sum_{i=1}^{p} \tau_{i} f_{i}(E(x)) + \sum_{t=1}^{m} \xi_{t} g_{t}(E(x)) < \sum_{i=1}^{p} \tau_{i} f_{i}(E(y)) + \sum_{t=1}^{m} \xi_{t} g_{t}(E(y)),$$
(11)

$$\sum_{i=1}^{p} \tau_{i} f_{i}(E(x)) < \sum_{i=1}^{p} \tau_{i} f_{i}(E(y)),$$
(12)

$$\sum_{t=1}^{m} \xi_{t} g_{t}(E(x)) \leq \sum_{t=1}^{m} \xi_{t} g_{t}(E(y)).$$
(13)

By assumption,  $x \in D_E$  and  $(y, \tau, \xi) \in \Gamma_E$ . Since (f, g) is *E*-type I with respect to  $\eta$  at y and by Definition 2.8, the following inequalities

$$f_i(E(x)) - f_i(E(y)) \ge \nabla f_i(E(y))\eta(E(x), E(y)), \quad i \in I$$
(14)

$$0 = -g_t(E(y)) \ge \nabla g_t(E(y))\eta(E(x), E(y)), \quad t \in T(E(y))$$
(15)

hold, respectively. By  $\tau_i > 0$ , i = 1, 2, ..., p,  $\xi_i \ge 0$ ,  $t \in T(E(y))$ , using above inequalities, we obtain that the inequality

$$\sum_{i=1}^{p} \tau_{i} f_{i}(E(x)) - \sum_{i=1}^{p} \tau_{i} f_{i}(E(y)) \geq \left[ \sum_{i=1}^{p} \tau_{i} \nabla f_{i}(E(y)) + \sum_{t \in T(E(y))} \xi_{t} \nabla g_{t}(E(y)) \right] \eta(E(x), E(y)) \quad (16)$$

holds. Thus, by the first constraint of  $(WD_E)$ , it follows that the following inequality

$$\sum_{i=1}^{p} \tau_{i} f_{i}(E(x)) \ge \sum_{i=1}^{p} \tau_{i} f_{i}(E(y))$$
(17)

holds, contradicting (12). Thus, the proof of this theorem is completed.  $\Box$ 

**Theorem 3.2 :** (Wolfe weak E-duality between (VP) and  $(WD_E)$ ). Let  $E(x) \in D$  and  $(y, \tau, \xi) \in \Gamma_E$ . Further, suppose that all hypotheses of Theorem 3.1 are fulfilled. Then, Wolfe weak E-duality between (VP) and  $(WD_E)$  holds, that is,

$$f(E(x)) \not\prec f(E(y)) + \left\lfloor \sum_{t=1}^{m} \xi_t g_t(E(y)) \right\rfloor e.$$

*Proof* : Let *E*(*x*) ∈ *D* and  $(y, τ, ξ) ∈ Γ_E$ . Then, Lemma 2.5 implies that  $x ∈ D_E$ . Since all hypotheses of Theorem 3.1 are fulfilled, from Theorem 3.1 the Wolfe weak *E*-duality between (VP) and (WD<sub>E</sub>) follows directly. □

**Theorem 3.3 :** (Wolfe weak duality between  $(VP_E)$  and  $(WD_E)$ ). Let  $x \in D_E$  and  $(y, \tau, \xi) \in \Gamma_E$  such that (f, g) is strictly *E*-type *I* at *y*. Then

$$f(E(x)) \not\prec f(E(y)) + \left[\sum_{t=1}^{m} \xi_t g_t(E(y))\right] e.$$
(18)

**Theorem 3.4 :** (Wolfe weak E-duality between (VP) and  $(WD_E)$ ). Let  $E(x) \in D$ and  $(y, \tau, \xi) \in \Gamma_E$ . Further, suppose that all hypotheses of Theorem 3.3 are fulfilled. Then, weak E-duality between (VP) and  $(WD_F)$  holds, that is,

$$f(E(x)) \not\prec f(E(y)) + \left[\sum_{t=1}^{m} \xi_t g_t(E(y))\right] e.$$

**Theorem 3.5 :** (Wolfe strong duality between  $(VP_E)$  and  $(WD_E)$  and also strong *E*-duality between (VP) and  $(WD_E)$ ). Let  $\overline{x} \in D_E$  be a weakly efficient solution (an efficient solution) of the problem  $(VP_E)$  (and, thus,  $E(\overline{x}) \in D$  be a weakly *E*-efficient solution (an *E*-efficient solution) of the problem (VP)). Further, assume that the Guignard constraint qualification [3] be satisfied at  $\overline{x}$ . Then there exist  $\tau \in R^p$ ,  $\xi \in R^m$ ,  $\xi \ge 0$  such that  $(\overline{x}, \tau, \xi) \in \Gamma_E$ . If all hypotheses of (Theorem 3.1) Theorem 3.3 are satisfied, then  $(\overline{x}, \tau, \xi)$  is a weakly efficient solution (an efficient solution) of a maximum type in the problem  $(WD_E)$ .

In other words, if  $E(\overline{x}) \in D$  is a weakly *E*-efficient solution (an *E*-efficient solution) of the problem (VP), then  $(\overline{x}, \tau, \xi)$  is a weakly efficient solution (an efficient solution) of a maximum type in the dual problem (WD<sub>*E*</sub>).

**Proof** : Since  $\overline{x} \in D_E$  is a weakly efficient solution of the problem  $(VP_E)$  and the Guignard constraint qualification [3] is satisfied at  $\overline{x}$ , there exist  $\tau \in R^p$ ,  $\xi \in R^m$ ,  $\xi \ge 0$  such that

$$\sum_{i=1}^{p} \tau_{i} \nabla (f_{i} \circ E)(\overline{x}) + \sum_{t=1}^{m} \xi_{t} \nabla (g_{t} \circ E)(\overline{x}) = 0,$$
$$\xi_{t} (g_{t} \circ E)(\overline{x}) = 0, t \in T,$$
$$\tau \ge 0, \xi \ge 0.$$

Thus,  $(\bar{x}, \tau, \xi)$  is a feasible solution for  $(WD_E)$ . If  $(\bar{x}, \tau, \xi)$  is not a weakly efficient solution (an efficient solution) for  $(WD_E)$ , then there exists a feasible solution  $(\tilde{x}, \tilde{\tau}, \tilde{\xi})$  of  $(WD_E)$  such that  $f(E(\tilde{x})) < f(E(\bar{x}))$ , which contradicts the Theorem 3.1. Hence  $(\bar{x}, \tau, \xi)$  is a weakly efficient solution (an efficient solution) for  $(WD_E)$ .

In addition, according to Lemma 2.5, we have  $E(\overline{x}) \in D$ . Since  $\overline{x} \in D_E$ is an efficient solution of the problem  $(VP_E)$ , Lemma 2.7 implies that  $E(\overline{x})$ is a weakly *E*-efficient solution in the problem (VP). Then, through the strong duality between  $(VP_E)$  and  $(WD_E)$ , we conclude that the strong *E*-duality also holds between the problems (VP) and  $(WD_E)$ . This implies that if  $E(\overline{x}) \in D$  is a weakly *E*-efficient solution of the problem (VP), there exist  $\tau \in R^P$ ,  $\xi \in R^m$ ,  $\xi \ge 0$  such that  $(\overline{x}, \tau, \xi)$  is a weakly efficient solution of a maximum type in the problem  $(WD_E)$ .

### 4. Mixed E-duality

In this section, for the considered *E*-differentiable multiobjective programming problem (VP), we define the following dual problem in the sense of mixed. Further, under *E*-type I functions, we prove *E*-duality results between the multiobjective programming problem (VP) and its vector *E*-dual problem in the sense of mixed.

Let the index set *T* be divided into two different subsets,  $T_1$  and  $T_2$ , so that  $T = T_1 \cup T_2$ . Let  $T_1$  be a subset of *T* and  $T_2 = T \setminus T_1$ ,  $\xi_{T_K} g_{T_K} = \sum_{t \in T_k} \xi_t g_t$ , k = 1, 2 and, in addition, the cardinality of index sets  $T_1$  and  $T_2$ , respectively, is denoted by  $|T_1|$  and  $|T_2|$ . Let

$$D^{2} = \{x \in R^{n} : g_{t}(x) \leq 0, t \in T_{2}\}.$$

Now, for the problem (VP), we introduce the definition of the scalar Lagrange function  $L: D \times R^p_+ \times R^{|T_1|}_+ \to R$  as follows

$$L(x,\tau,\xi_{T_1}) := \sum_{i=1}^p \tau_i f_i(x) + \sum_{t \in T_1} \xi_t g_t(x).$$
(19)

Further, let  $E: \mathbb{R}^n \to \mathbb{R}^n$  be a given one-to-one and onto operator. Furthermore, let

$$D_E^2 := \{x \in R^n : (g_t \circ E)(x) \leq 0, t \in T_2\}.$$

Now, we define for problem (VP) its *E*-differentiable vector *E*-dual problem in the sense of mixed as follows:

maximize 
$$f(E(y)) + \left[ \xi_{T_1}(g_{T_1} \circ E)(y) \right] e$$
  
subject to  $\tau \nabla f(E(y)) + \xi \nabla g(E(y)) = 0$ , (MVD<sub>E</sub>)  
 $\xi_{T_2}(g_{T_2} \circ E)(y) \ge 0$ ,  
 $\tau \in \mathbb{R}^p, \tau \ge 0, \tau e = 1, \xi \in \mathbb{R}^m, \xi \ge 0.$ 

Let  $\Psi_E$  denote the set of all feasible solutions of  $(MVD_E)$ , that is,  $\Psi_E = \left\{ (y, \tau, \xi) \in \mathbb{R}^n \times \mathbb{R}^p \times \mathbb{R}^q : \tau \nabla f(E(y)) + \xi \nabla g(E(y)) = 0, \\ \sum_{t \in T_2} \xi_t(g_t \circ E)(y) \ge 0, \tau \ge 0, \tau e = 1, \xi \ge 0 \right\}.$  Further, let  $Y_E = \{y \in \mathbb{R}^n : (y, \tau, \xi) \in \Psi_E\}$ . Note that we get a vector Wolfe *E*-dual problem [4] for  $T_1 = \emptyset$  and a vector Mond-Weir *E*-dual problem [1], [5] for  $T_2 = \emptyset$  in the vector mixed *E*-dual problem (MVD<sub>E</sub>), respectively.

**Theorem 4.1**: (*Mixed weak duality between*  $(VP_E)$  and  $(MVD_E)$ ). Let  $x \in D_E$  and  $(y, \tau, \xi) \in \Psi_E$  such that (f, g) is E-type I at y. Then

$$f(E(x)) \not< f(E(y)) + [\xi_{T_1}(g_{T_1} \circ E)(y)]e.$$
 (20)

*Proof* : Let *x* ∈ *D*<sub>*E*</sub> and (*y*, *τ*, *ξ*) ∈  $\Psi_E$ . By means of contradiction, suppose that

$$(f \circ E)(x) < (f \circ E)(y) + \left[\xi_{T_1}(g_{T_1} \circ E)(y)\right]e.$$

Thus,

$$(f_i \circ E)(x) < (f_i \circ E)(y) + \left[\sum_{t \in T_1} \xi_t(g_t \circ E)(y)\right], i \in I.$$

$$(21)$$

Multiplying (21) by  $\tau_i$  and then adding both sides of the resulting inequalities, we get

$$\sum_{i=1}^{p} \tau_{i}(f_{i} \circ E)(x) < \sum_{i=1}^{p} \tau_{i}(f_{i} \circ E)(y) + \left[\sum_{t \in T_{1}} \xi_{i}(g_{t} \circ E)(y)\right] \sum_{i=1}^{p} \tau_{i}.$$

Since  $\sum_{i=1}^{p} \tau_i = 1$ , the following inequality

$$\sum_{i=1}^{p} \tau_{i}(f_{i} \circ E)(x) < \sum_{i=1}^{p} \tau_{i}(f_{i} \circ E)(y) + \sum_{t \in T_{1}} \xi_{t}(g_{t} \circ E)(y)$$

holds. By  $x \in D_E$  and  $(y, \tau, \xi) \in \Psi_E$ , we have

$$\sum_{i=1}^{p} \tau_{i}(f_{i} \circ E)(x) + \sum_{t \in T_{1}} \xi_{t}(g_{t} \circ E)(x) < \sum_{i=1}^{p} \tau_{i}(f_{i} \circ E)(y) + \sum_{t \in T_{1}} \xi_{t}(g_{t} \circ E)(y),$$
(22)

$$\sum_{i=1}^{p} \tau_{i}(f_{i} \circ E)(x) < \sum_{i=1}^{p} \tau_{i}(f_{i} \circ E)(y),$$
(23)

$$\sum_{t\in T_2} \xi_t(g_t \circ E)(x) \leq \sum_{t\in T_2} \xi_t(g_t \circ E)(y),$$
(24)

Since (f,g) is *E*-type I with respect to  $\eta$  at *y* and by Definition 2.8, the following inequalities

$$f_i(E(x)) - f_i(E(y)) \ge \nabla f_i(E(y))\eta(E(x), E(y)), \quad i \in I$$
(25)

$$0 = -g_t(E(y)) \ge \nabla g_t(E(y))\eta(E(x), E(y)), \quad t \in T(E(y))$$
(26)

hold, respectively. By  $\tau_i > 0$ , i = 1, 2, ..., p,  $\xi_t \ge 0$ ,  $t \in T(E(y))$ , using above inequalities, we obtain that the inequality

$$\sum_{i=1}^{p} \tau_{i} f_{i}(E(x)) - \sum_{i=1}^{p} \tau_{i} f_{i}(E(y)) \ge \left[ \sum_{i=1}^{p} \tau_{i} \nabla f_{i}(E(y)) + \sum_{t \in T(E(y))} \xi_{t} \nabla g_{t}(E(y)) \right] \eta(E(x), E(y))$$
(27)

holds. Thus, by the first constraint of  $(MVD_E)$ , it follows that the following inequality

$$\sum_{i=1}^{p} \tau_{i} f_{i}(E(x)) \ge \sum_{i=1}^{p} \tau_{i} f_{i}(E(y))$$
(28)

holds, contradicting (23). Thus, the proof of this theorem is completed.

**Theorem 4.2**: (*Mixed weak E-duality between* (*VP*) and (*MVD<sub>E</sub>*)). Let  $E(x) \in D$ and  $(y, \tau, \xi) \in \Psi_E$ . In addition, assume that all hypotheses of Theorem 4.1 are satisfied. Then, mixed weak E-duality between (*VP*) and (*MVD<sub>E</sub>*) holds, that is,

$$f(E(x)) \not< f(E(y)) + \left[ \xi_{T_1}(g_{T_1} \circ E)(y) \right] e.$$

**Theorem 4.3 :** (*Mixed strong duality between*  $(VP_E)$  *and*  $(MVD_E)$  *and also strong E*-duality between (VP) and  $(MVD_E)$ ). Let  $\overline{x} \in D_E$  be a weakly efficient solution (an efficient solution) of the problem  $(VP_E)$  (and, thus,  $E(\overline{x}) \in D$  be a weakly *E*-efficient solution (an *E*-efficient solution) of the problem (VP)). Further, assume that the Guignard constraint qualification [3] be satisfied at  $\overline{x}$ . Then there exist  $\tau \in \mathbb{R}^p$ ,  $\xi \in \mathbb{R}^m$ ,  $\xi \ge 0$  such that  $(\overline{x}, \tau, \xi) \in \Psi_E$ . If all hypotheses of (Theorem 4.1) are satisfied, then  $(\overline{x}, \tau, \xi)$  is a weakly efficient solution (an efficient solution) of a maximum type in the problem (MVD\_E).

In other words, if  $E(\overline{x}) \in D$  is a weakly *E*-efficient solution (an *E*-efficient solution) of the problem (VP), then  $(\overline{x}, \tau, \xi)$  is a weakly efficient solution (an efficient solution) of a maximum type in the dual problem (MVD <sub>*E*</sub>).

### 5. Concluding remarks

In this paper, new classes of *E*-differentiable multiobjective programming problems have been considered. For such *E*-differentiable multiobjective programming problems, the so-called vector Wolfe and mixed *E*-dual problems have been formulated. Then, under *E*-type I functions, various *E*-duality theorems have been proven between the considered *E*-differentiable vector optimization problem and its vector Wolfe and mixed *E*-duals. Vector Wolfe type and vector Mond-Weir type *E*-duals are special cases of a vector mixed type *E*-dual.

# References

- Antczak, T., and N. Abdulaleem. "E-optimality conditions and Wolfe E-duality for E-differentiable vector optimization problems with inequality and equality constraints." *Journal of Nonlinear Sciences and Applications*. 12, 745--764 (2019).
- [2] Antczak, T., and N. Abdulaleem. "Optimality and duality results for E-differentiable multiobjective fractional programming problems under E-convexity." *Journal of Inequalities and Applications*. 2019, 292 (2019).
- [3] Abdulaleem, N. "E-invexity and generalized E-invexity in E-differentiable multiobjective programming." ITM Web of Conferences. Vol. 24. EDP Sciences, (2019).
- [4] Abdulaleem, N. "E-duality results for E-differentiable E-invex multiobjective programming problems." *Journal of Physics: Conference Series*. (Vol. 1294, No. 3, p. 032027). IOP Publishing, (2019).
- [5] Abdulaleem, N. "Wolfe E-duality for E-differentiable E-invex vector optimization problems with inequality and equality constraints." International Conference on Applied Mathematics & Computational Science (ICAMCS.NET). IEEE, Budapest, Hungary, 156-163 (2018).
- [6] Abdulaleem, N. "E-optimality conditions for E-differentiable E-invex multiobjective programming problems." WSEAS Transactions on Mathematics. 18, 14-27 (2019).
- [7] Abdulaleem, N. "Optimality and duality for *E*-differentiable multiobjective programming problems involving *E*-type I functions." submitted for publication, 2020.

- [8] Aghezzaf, B. and Hachimi, M. "Generalized invexity and duality in multiobjective programming problems." *Journal of Global Optimization*. 18, 91-101 (2000).
- [9] Hanson, M. A. "On sufficiency of the Kuhn-Tucker conditions." *Journal of Mathematical Analysis and Applications*. 80, 545-550 1981).
- [10] Hanson, M.A. and Mond, B. "Necessary and sufficient conditions in constrained optimization." *Mathematical programming*. 37, 51-58 (1987).
- [11] Kaul, R. N., Suneja, S. K. and Srivastava, M. K. "Optimality criteria and duality in multiple-objective optimization involving generalized invexity." *Journal of Optimization Theory and Applications*. 80, 465-482 (1994).
- [12] Megahed, Abd El-Monem A., Gomma, H. G., Youness, E. A. and El-Banna, A. Z. "Optimality conditions of *E*-convex programming for an *E*-differentiable function." *Journal of Inequalities and Applications*. 2013, 246 (2013).
- [13] Rueda, N. G. and Hanson, M.A. "Optimality criteria in mathematical programming involving generalized invexity." *Journal of Mathematical Analysis and Applications*. 130, 375-385 (1988).
- [14] Xu Z. "Mixed type duality in multiobjective programming problems." *Journal of Mathematical Analysis and Applications*, 198(3), 621-635 (1996).
- [15] Youness, E. A. "E-convex sets, E-convex functions, and E-convex programming." Journal of Optimization Theory and Applications. 102, 439-450 (1999).

Received June, 2020 Revised November, 2020